Binary trees Basics of Programming 1



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Content



Binary trees

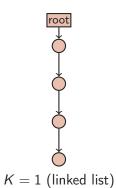
Chapter 1

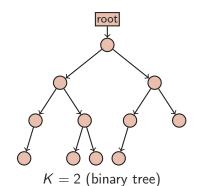
Binary trees



Trees



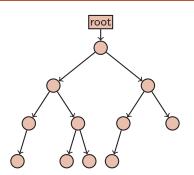




- An acyclic graph
- Every node has exactly one incoming edge
- K-ary tree: every node has at most K outgoing edges

Binary trees





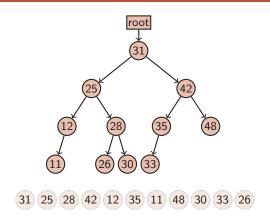
Declaration of the binary tree data structure

```
typedef struct tree {
int data;
struct tree *left, *right;
tree_elem, *tree_ptr;
ink
```

■ Typically we typedef not only the struct, but also the pointer

Binary search trees

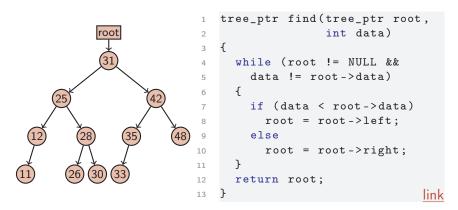




- Sub-tree to the left: only elements smaller than the node
- Sub-tree to the right: only elements greater than the node
- The structure of the tree depends on the insertion order of the elements!

Searching an element in the tree

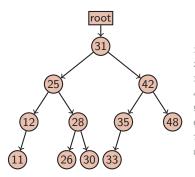




- This is not recursive yet
- In a depth-d tree the max. number of steps is d
- If the tree is balanced and has n elements $\Rightarrow \approx \log_2 n$ steps!

In-order traversal





```
void inorder(tree_ptr root)
  if (root == NULL)
    return:
  inorder(root->left);
  printf("%d ", root->data);
  inorder(root->right);
```

11 12 25 26 28 30 31 33 35 42 48

- in-order traversal
 - left sub-tree
 - root element
 - right sub-tree

With this traversal the nodes are visited in increasing order of their values

In-order traversal



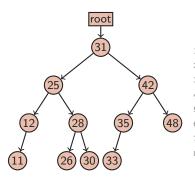
An other implementation of the in-order traversal:

```
void inorder(tree_ptr root)
2
    if (root->left != NULL)
       inorder(root->left);
    printf("%d ", root->data);
    if (root->right != NULL)
6
       inorder(root->right);
  }
8
```

But in this case the caller has not make sure that root != NULL holds

Pre-order traversal





```
void preorder(tree_ptr root)
 if (root == NULL)
    return;
  printf("%d ", root->data);
  preorder(root->left);
  preorder(root->right);
```

31 25 12 11 28 26 30 42 35 33 48

- pre-order traversal
 - root element
 - left sub-tree
 - right sub-tree

Saving the elements of the tree in this order, and building it again, the structure of the tree can be fully reconstructed.

Building a tree



Inserting a new node to the tree

```
tree_ptr insert(tree_ptr root, int data)
2
     if (root == NULL) {
3
       root = (tree_ptr)malloc(sizeof(tree_elem));
4
       root -> data = data;
5
     }
6
    else if (data < root->data)
7
       root->left = insert(root->left, data);
8
    else
9
       root -> right = insert(root -> right, data);
10
     return root;
11
                                                           link
12
```

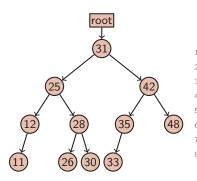
Usage of this function:

```
tree_ptr root = NULL;
root = insert(root, 2);
  root = insert(root, 8);
  . . .
```

Binary trees

Post-order traversal





```
void postorder(tree_ptr root)
  if (root == NULL)
    return:
  postorder(root->left);
  postorder(root->right);
  printf("%d ", root->data);
```

12 26 30 28 25 33 35 48 42 31

- post-order traversal
 - left sub-tree
 - right sub-tree
 - root element

In this order the leaves of the tree are visited first \rightarrow application: releasing/deleting a tree

Deleting a tree by post-order traversal



```
void delete(tree_ptr root)
2
    if (root == NULL) /* empty tree: nothing to delete >
      return;
4
  delete(root->left); /* post-order traversal */
    delete(root->right);
    free(root);
7
                                                       link
```

A program segment (without memory leaks):

```
tree_ptr root = NULL;
root = insert(root, 2);
  root = insert(root, 8);
  . . .
5 delete(root);
6 root = NULL;
```

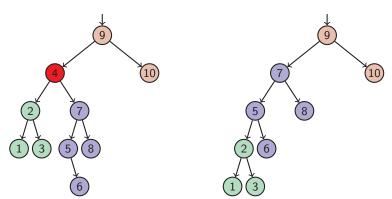
Simple algorithms on binary trees



- Write a recursive function (max. 10 lines), that
 - determines the depth of a tree
 - calculates the count / the sum / the average of the values stored in the nodes of the tree
- Write a iterative function (max. 10 lines), that
 - computes the minimum and the maximum of the values stored in the nodes
 - returns the pointer to the node storing the maximal / minimal value of the tree

Deleting an element from a search tree – naively

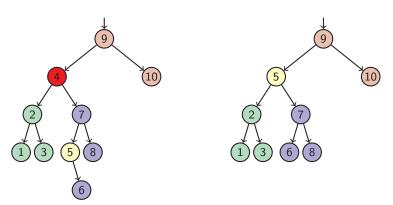




- The right sub-tree is moved to the place of the deleted node
- The left sub-tree is inserted to below the minimal element of the right sub-tree
- The tree is getting imbalanced!

Deleting an element from a search tree – clever

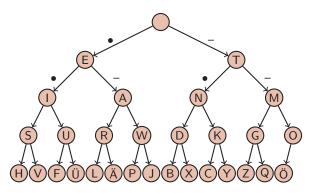




- The minimal element of the right sub-tree is moved to the place of the deleted node
- This element could have only a right sub-tree, it is moved one level up, to its old place

Morse decoding tree



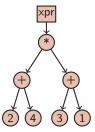


SOSOS:

The response:

Evaluating mathematical expressions





- Storing math expressions in a tree
- Leaves → numeric constants
- Branches → two-operand operators
- In the example: (2+4)*(3+1)

```
int eval(tree_ptr xpr)
2
    char c = xpr->data;
3
    if (isdigit(c)) /* stopping condition */
      return c - '0';
5
    if (c == '+')
      return eval(xpr->left) + eval(xpr->right);
    if (c == '*')
      return eval(xpr->left) * eval(xpr->right);
9
```

Evaluating functions



Let us introduce variable x as a leaf node as well:

```
double feval(tree_ptr xpr, double x)
2
     char c = xpr->data;
3
     if (isdigit(c))
       return c - '0';
5
     if (c == 'x')
6
       return x;
     if (c == '+')
8
       return feval(xpr->left, x) + feval(xpr->right, x);
9
     if (c == '*')
10
       return feval(xpr->left, x) * feval(xpr->right, x);
11
   }
                                                           link
12
```

Evaluating the derivative of a function



Let us take the derivative of the function! The rules are:

```
• c' = 0

• x' = 1

• (f + g)' = f' + g'

• (f \cdot g)' = f' \cdot g + f \cdot g'
```

```
double deval(tree_ptr xpr, double x)
2
3
     char c = xpr->data;
     if (isdigit(c)) /* stopping condition */
       return 0.0;
5
   if (c == 'x')
                       /* stopping condition */
6
       return 1.0:
7
   if (c == '+')
       return deval(xpr->left, x) + deval(xpr->right, x);
9
     if (c == '*')
10
       return deval(xpr->left, x) * feval(xpr->right, x) +
11
         feval(xpr->left, x) * deval(xpr->right, x);
12
                                                        link
13
```

Thank you for your attention.